

# Treeningvõistlus

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1. (a) Leia kõik sellised pidevad funktsioonid  $f : [0, \infty) \rightarrow \mathbb{R}$ , et iga  $x \geq 0$  korral

$$f(x^2) = f(x).$$

- (b) Kas leidub lõpmata palju erinevaid pidevaid funktsioone  $f : (1, \infty) \rightarrow \mathbb{R}$ , mille jaoks iga  $x > 1$  korral

$$f(x^2) - f(x) = \ln 2?$$

2. Olgu  $A$  ja  $B$  reaalsed  $n \times n$  maatriksid. Tõesta, et

$$A(A + B)^{-1}B = B(A + B)^{-1}A$$

eeldusel, et

- (a) maatriksid  $A + B$ ,  $A$  ja  $B$  on pööratavad,  
(b) maatriks  $A + B$  on pööratav.

3. Iga  $n \in \mathbb{N}$  korral tähistame

$$S_n = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \cdots + \frac{1}{2n-1} - \frac{1}{2n}.$$

Tõesta, et

$$S_n = \frac{1}{n+1} + \frac{1}{n+2} + \cdots + \frac{1}{2n}$$

ja leia

$$\lim_{n \rightarrow \infty} S_n.$$

4. Milliste  $m \geq 1$  korral leiduvad täisarvud  $n_0 < n_1 < \dots < n_m$  ja täisarvuliste kordajatega polünoom  $P$ , milleaste on  $m$ , nii, et

$$|P(n_0)| = |P(n_1)| = \dots = |P(n_m)| = 1?$$

## Solutions

1. (a) These are constant functions. They clearly satisfy the condition. On the other hand, fix  $x > 0$  and denote  $x_n := x^{\frac{1}{2^n}}$ . Then  $x_n \rightarrow 1$ . So  $f(x_n) \rightarrow f(1)$ . But

$$f(x) = f(x_0) = f(x_1) = \dots = f(x_n),$$

so  $f(x) = f(1)$ . Then, by continuity, also  $f(0) = f(1)$ .

- (b) Yes, because  $h(x) = \ln \ln x$  as well as  $h(x) + C$ , where  $C \in \mathbb{R}$  is a constant, satisfy the condition.

2. (a) We have

$$\begin{aligned} A(A+B)^{-1}B &= (A^{-1})^{-1}(A+B)^{-1}(B^{-1})^{-1} \\ &= (B^{-1}(A+B)A^{-1})^{-1} \\ &= (B^{-1} + A^{-1})^{-1}. \end{aligned}$$

Since the latter expression is symmetric with respect to  $A$  and  $B$ , we get the claim.

- (b) This is because

$$A(A+B)^{-1}B + B(A+B)^{-1}B = (A+B)(A+B)^{-1}B = B$$

and

$$B(A+B)^{-1}A + B(A+B)^{-1}B = B(A+B)^{-1}(A+B) = B.$$

3. Set

$$H_n := 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n}$$

Then

$$S_n = H_{2n} - 2\left(\frac{1}{2}H_n\right) = H_{2n} - H_n.$$

To find a limit, observe that

$$\lim_n S_n = \lim_n \frac{1}{n} \sum_{i=1}^n \frac{1}{1 + \frac{i}{n}} = \int_0^1 \frac{dx}{1+x} = \ln 2.$$

4. For  $m = 1$ :  $P(x) = x$ ,  $\{n_i\} = \{-1, 1\}$ .

For  $m = 2$ :  $P(x) = x(x-1) - 1$ ,  $\{n_i\} = \{-1, 0, 1\}$ .

For  $m = 3$ :  $P(x) = (x+1)x(x-2) + 1$ ,  $\{n_i\} = \{-1, 0, 1, 2\}$ .

There are no such polynomials for  $m \geq 4$ . Using the Lagrange interpolation, we can show that

$$\max |P(n_i)| \geq |a_m| \frac{m!}{2^m} \geq \frac{m!}{2^m} \geq \frac{4!}{2^4} = \frac{24}{16} > 1$$

where  $a_m$  is the coefficient at  $x^m$ . Indeed, if we denote

$$f(x) = (x - n_0)(x - n_1)(x - n_2) \dots (x - n_m),$$

then the Lagrange interpolation formula gives

$$P(x) = \sum_{i=0}^m \frac{P(n_i)}{f'(n_i)} \frac{f(x)}{x - n_i}.$$

Looking at the coefficient at  $x_m$  we obtain

$$a_m = \sum_{i=0}^m \frac{P(n_i)}{f'(n_i)}$$

and

$$|a_m| \leq \max |P(n_i)| \sum_{i=0}^m \frac{1}{|f'(n_i)|}.$$

Note that

$$|f'(n_i)| = |(n_i - n_0)(n_i - n_1) \dots (n_i - n_{i-1})(n_i - n_{i+1}) \dots (n_i - n_m)| \geq i!(m-i)!.$$

Hence

$$\sum_{i=0}^m \frac{1}{|f'(n_i)|} \leq \frac{1}{m!} \sum_{i=0}^m \frac{m!}{i!(m-i)!} = \frac{2^m}{m!}.$$